

FIGURE 2

HIGHER LAYER	<u>240</u>
ATM ADAPTATION LAYER (AAL)	230
ATM LAYER	<u>220</u>
PHYSICAL LAYER	<u>210</u>

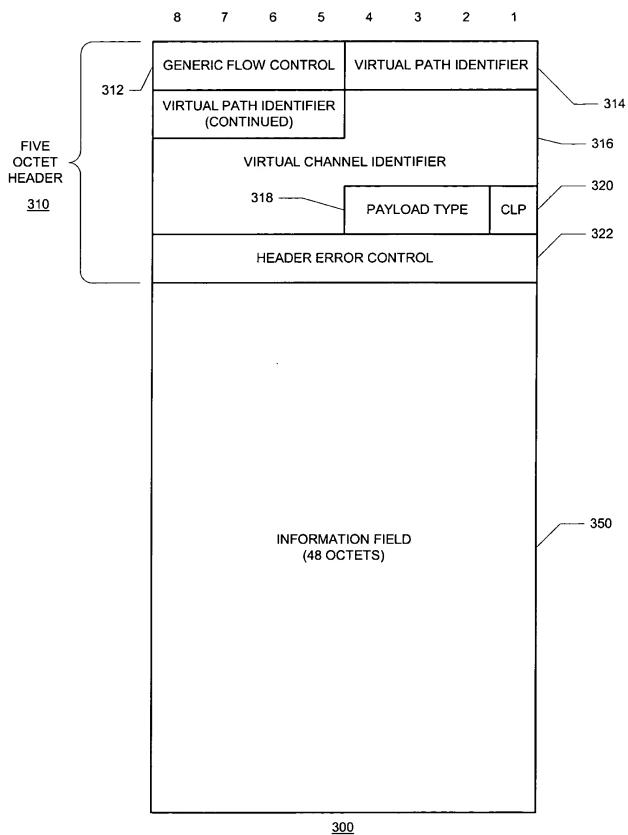
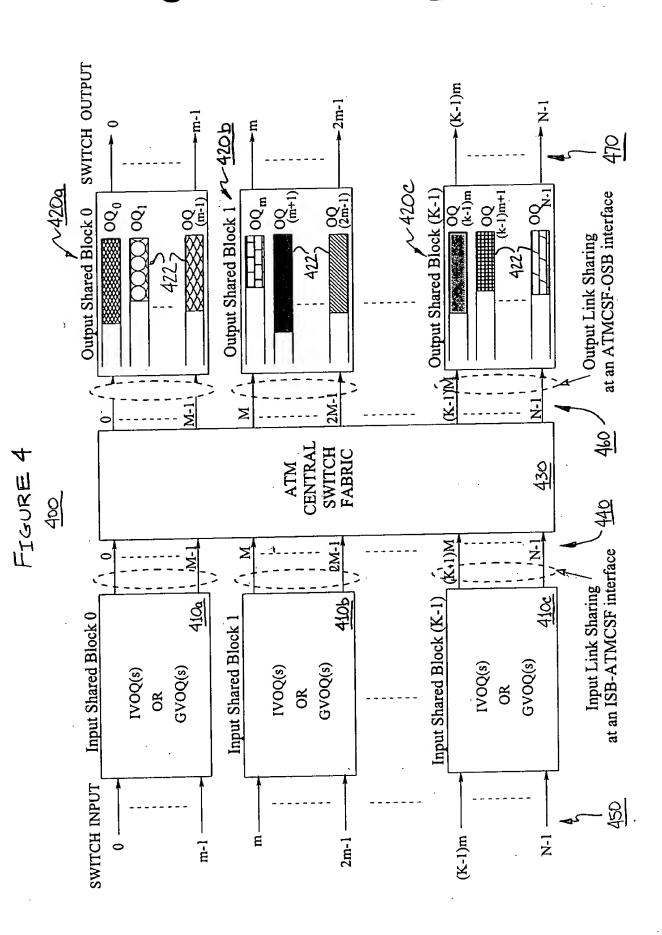


FIGURE 3A

FIGURE 3B



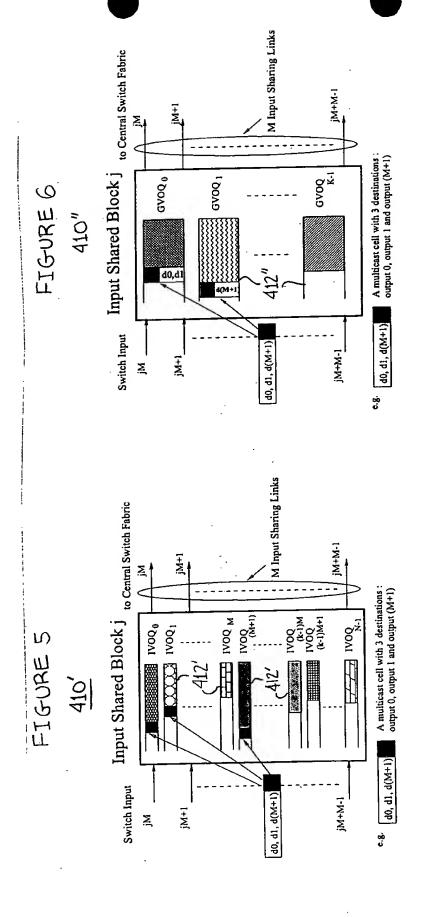
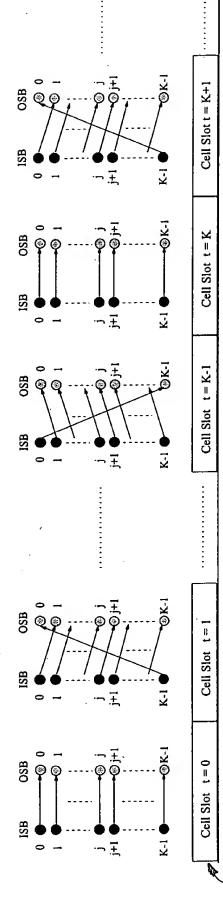
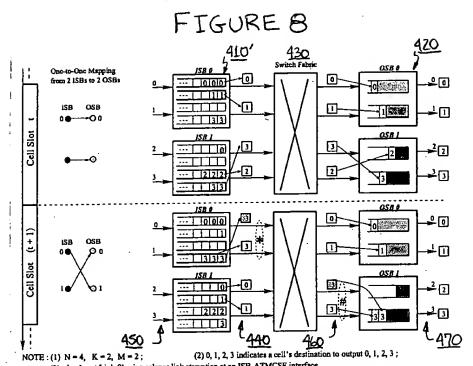


FIGURE 7

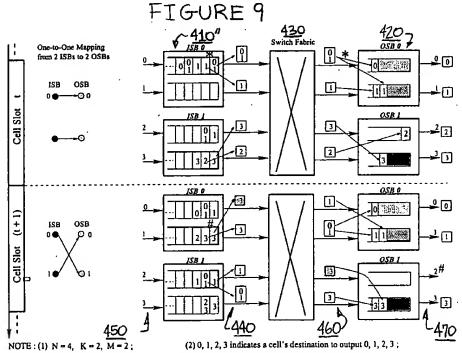


-710

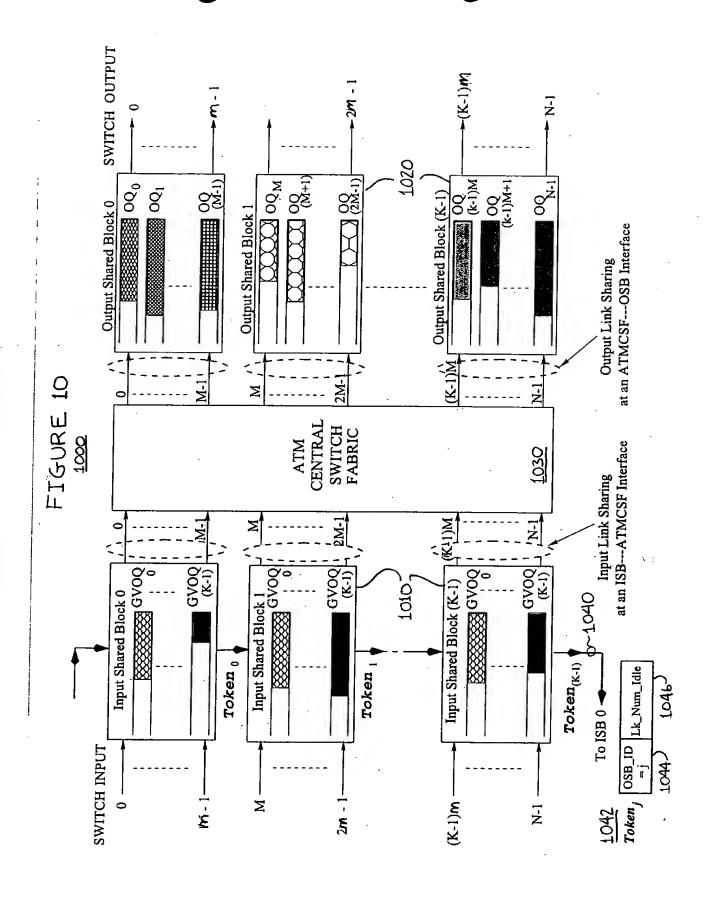


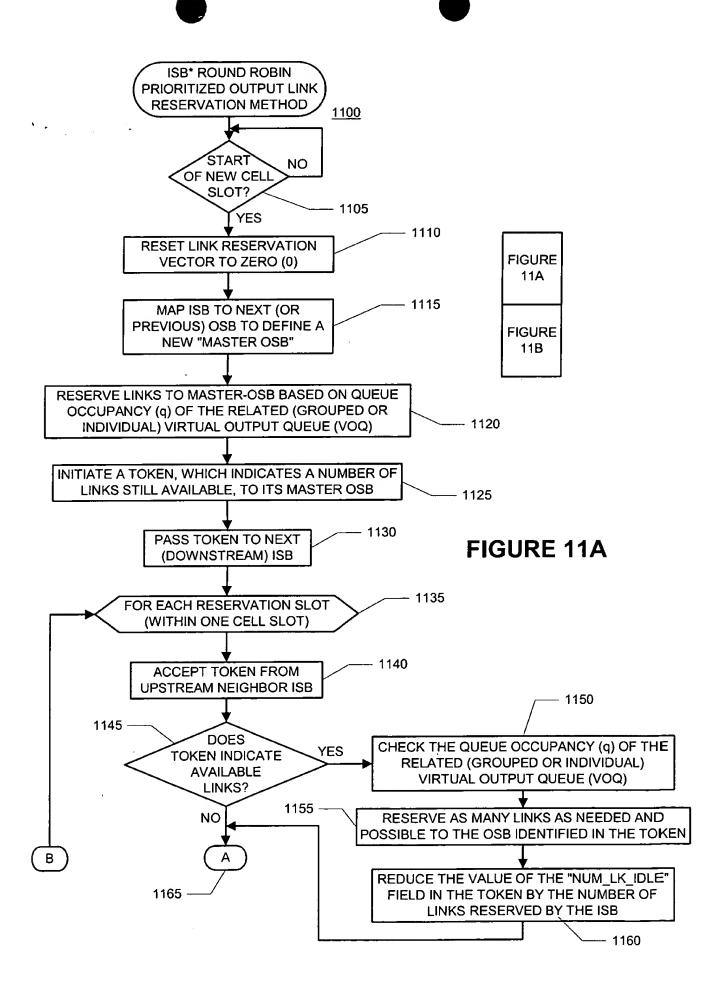
(3) * : Input Link Sharing relaxes link starvation at an ISB-ATMCSF interface.

(4) # : Output Link Sharing eliminates speedup requirement in central switch fabric; ATMCSF needs to keep cell sequence.



(3) * : Advantages of GVOQ, i.e. scheduling is simpler, many cell copies are forwarded.
 (4) # : Disadvantage of GVOQ, i.e. the cell going to output 2 is blocked by two cells destined to output 3.





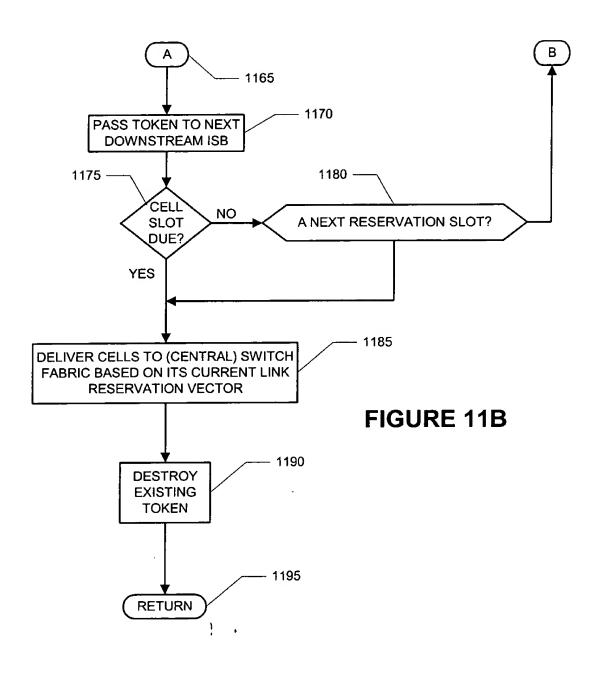
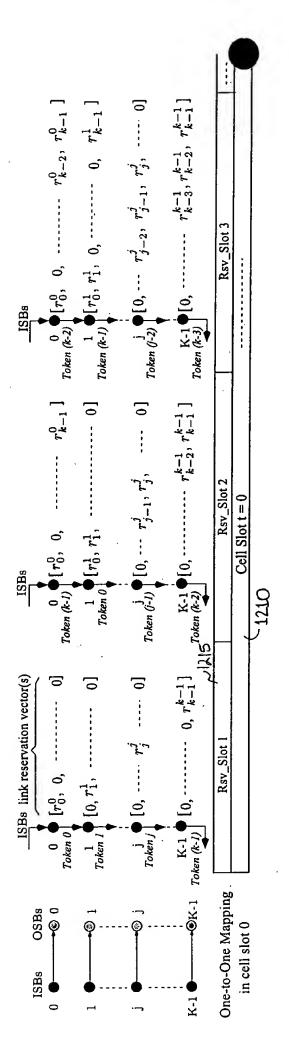
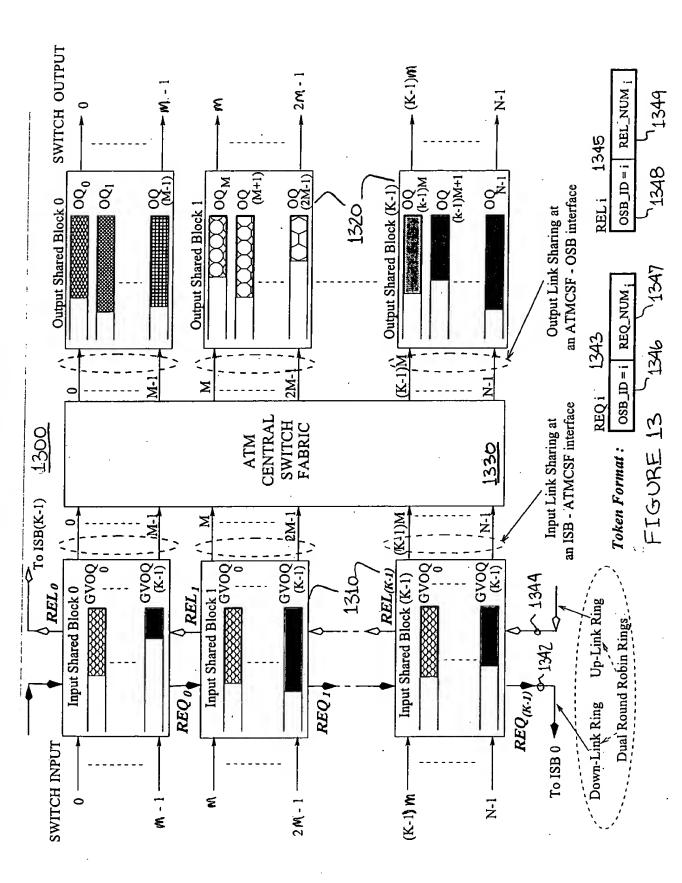


FIGURE 12





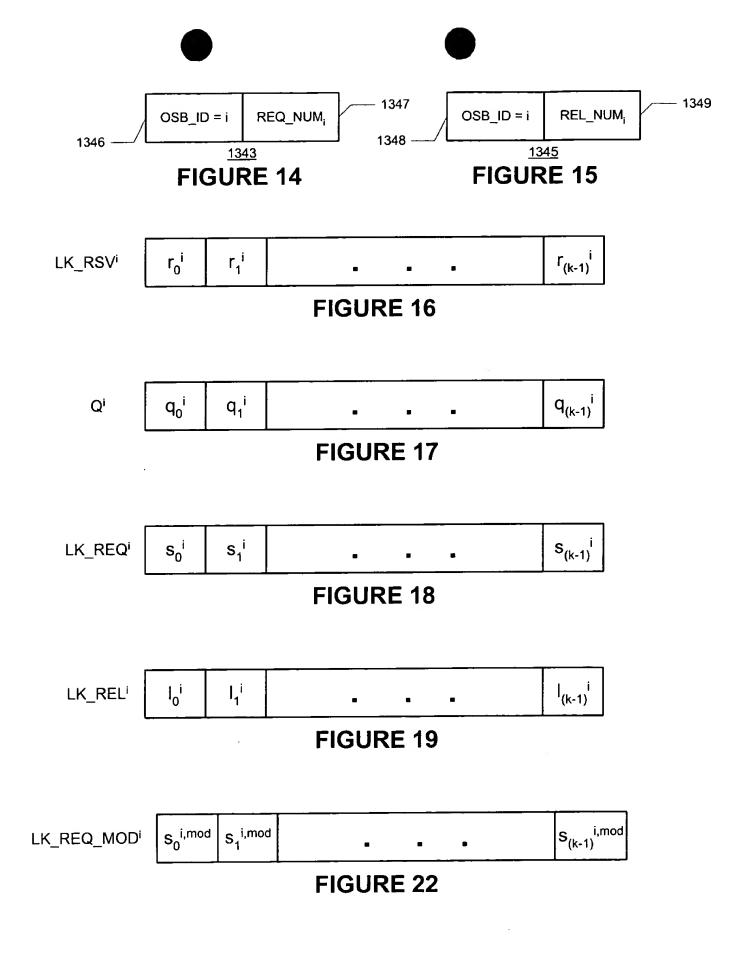


FIGURE 20

Upon Receiving REQ j the i^{th} ISB will request 0 extra link to the j^{th} OSB; the i^{tra} if ((REQ_NUM, > 0) and $(r_j^i > 0)$) if ((REQ_NUM, > 0) and then the ISB will release a link for OSB j; then the ISB will release if (REL_NUM, > 0), and ($\sum_{j=0}^{(h-1)} r_j^i < M$), if (REL_NUM, > 0), and ($\sum_{j=0}^{(h-1)} r_j^i < M$), then then	the j^{th} OSB; if ((REQ_NUM; > 0) and $(r_j^i) > Fair$) then the ISB will release a link for OSB; $< M$), a extra link to the n^{th} OSB)
the i th ISB take a link from REL _n token;	n token ;

REL n

FIGURE 21

	IF	7		THEN
	d_j^i	$S_j^{i,old}$	St	REQ.NUM;
case I	∴ in \	0	1	In Step 3, REQ_NUM, ++;
case 2		ı	1	REQ_NUM,
case 3		0	0	REQ_NUM,
case 4	≥ HT	.	if (REQ.) $s_j^i = 0$ else $s_i^i = 1$	N 25 1

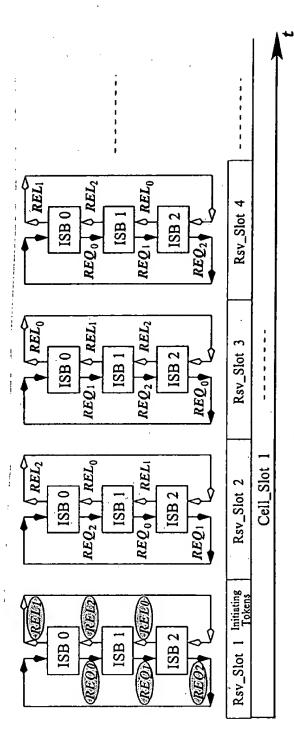
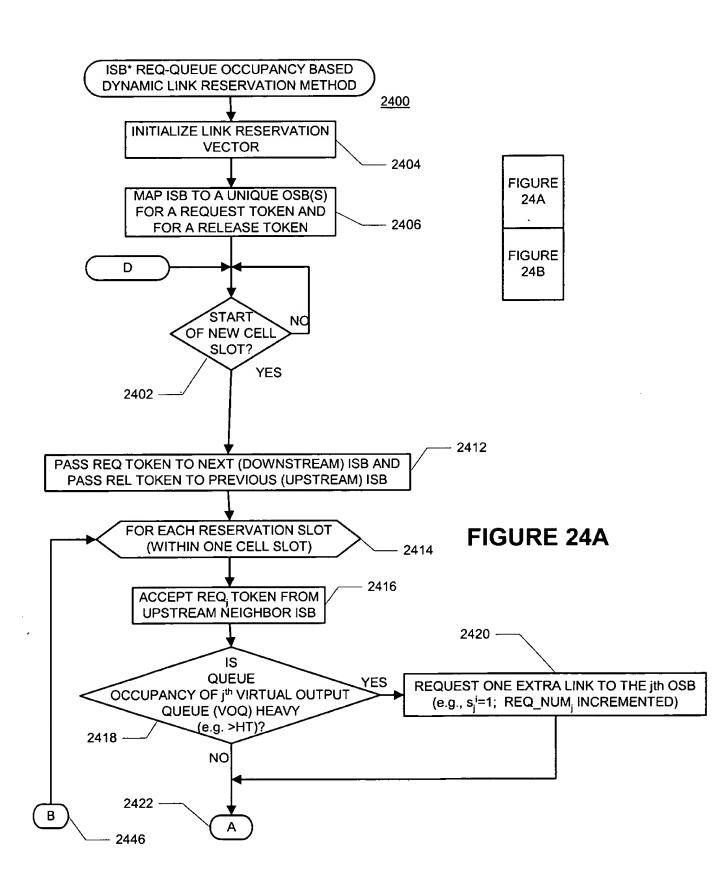
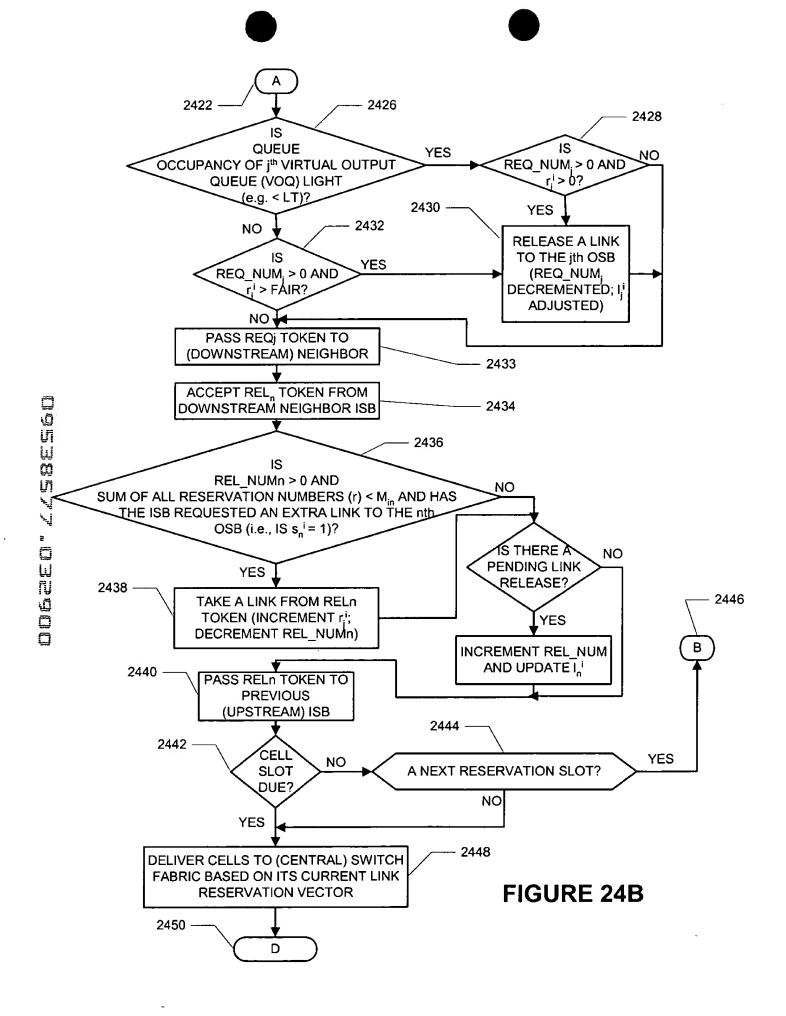


FIGURE 23





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FIGURE 25

Brown and		T	•						
Ouene Occupancy	the t th ISB will request 1 extra link to the j th OSB;	$(r_j^i > Fair))$	k for OSB j;	$Q_{m{q}} = \{ m{Q}_{m{q}} = \{ m{Q}_{m{q}} \} \} $ (Queue 'Occupancy, $\ m{Q}_{m{q}} \ _{m{Q}} \}$) where $\ m{Q}_{m{q}} \ _{m{Q}} = \{ m{Q}_{m{q}} \ _{m{Q}} \}$ (Leavy, load),	if (REL_NUM _n > 0)	and $(\sum_{l=0}^{(k-1)} r_l^i < M)$	then the ten	the t^{**} ISB will take a link from REL_{r} token:	•
Obele Occupancy $LT \leq q_1 \leq HT$ $LT \leq q_2 \leq HT$	ISB will request 0 extra link to the j^{th} OSB;	if $((REQ.NUM_j > 0)$ and $(r_j^i > Fair))$	then the ISB will release a link for OSB j;	$\ D_{t}^{(i)} \ = 0$ (leave Occupancy with $\ D_{t}^{(i)} \ \leq q_{t} \ \leq t H T$ (larea 2.5 normal loads)	if (REL_NUM, > 0)	and $(\sum_{i=0}^{(k-1)} r_i^i < M)$	and (the ISB has requested a link	to the $n^{""}$ USB)	the i^{th} ISB will take a link from REL_n token:
$(0) = \frac{(0) \cos (-(0) \cos (-(0)))}{(0)} = \frac{1}{2} \left[\frac{1}$	the i th ISB will request 0 ext	if $((REQ_NUM_j > 0) \text{ and } (r_j^i > 0))$	then the ISB will release a link for OSB j;	$q_{\pi}^{\mu} < \mu \mathcal{I}$ (Cueue: Occupancy (carea μ - Ilgh(load)).			then the ISB will release a link		
Operation Upon Receiving REQ j (QOBDLR)					Operations Upon	Receiving REL n	(KOBDEN)		
REQ.j	ISB i	←	REL n	HT LT	(6)		(1111)		

